

Towards Compositional Feedback in Non-Deterministic and Non-Input-Receptive Systems

Stavros Tripakis^{1,2}

Joint LICS'16 paper with Viorel Preteasa²

Contributions to the RCRS framework also by Iulia Dragomir²

¹UC Berkeley, USA

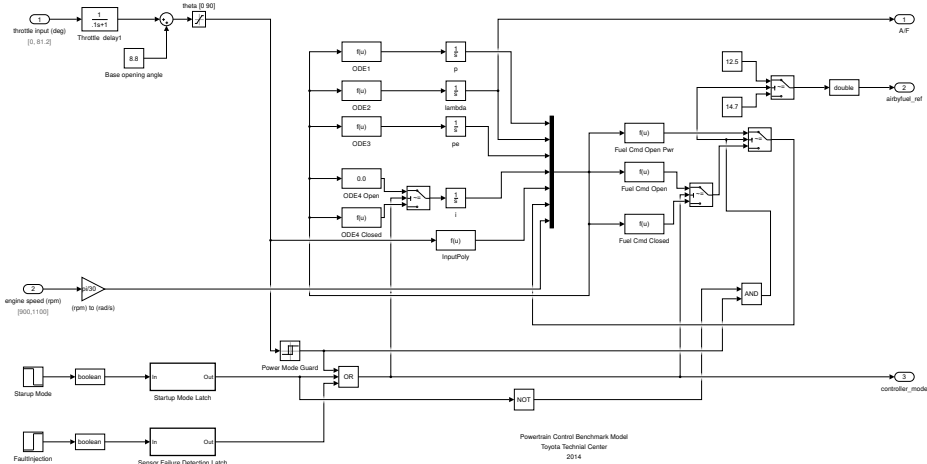
²Aalto University, Finland

ExCAPE PI Meeting 2016



Motivation: compositional reasoning for Simulink

Fuel Control System Model This model uses only the ODEs to implement the dynamics.



Benchmark provided by Toyota

Powertrain Control Benchmark Model
Toyota Technical Center
2014

This is a model of a hybrid automaton with polynomial dynamics, and an implementation of the 3rd model that appears in "Powertrain Control Verification Benchmark", 2014 Hybrid Systems: Computation and Control, X. Jin, J. V. Deshmukh, J. Kapinski, K. Ueda, and K. Butts

What does “compositional reasoning for Simulink” mean?

- 1 Be able to model basic Simulink blocks: constants, adders, integrators, ...
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- 4 Be able to **synthesize system from subsystems**: compute component compositions bottom-up.
- 5 Be able to **check substitutability**: component **refinement**.

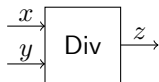
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- 6 Be able to express and verify **safety and liveness** properties.

Refinement Calculus of Reactive Systems (RCRS)

[EMSOFT 2009, TOPLAS 2011, SPIN 2013, EMSOFT 2014, FPS 2014, SPIN 2016, LICS 2016]

- Relational interfaces: symbolic, synchronous version of Alfaró-Henzinger's *interface automata*
- Example: relational interface for division component



contract: $y \neq 0 \wedge z = \frac{x}{y}$

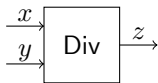
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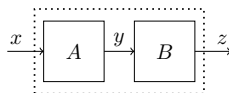
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- RCRS:

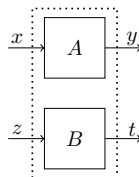
- Extends relational interfaces to handle liveness properties
- Semantics based on **monotonic property transformers**
- Can handle stateful systems, and **continuous** blocks by Euler discretization, e.g., Integrator: $s' = s + x \cdot dt$

Composition operators

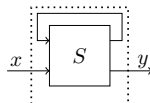
- Serial composition ($\forall\text{-}\exists$ synthesis inside!)



- Parallel composition (**conjunction**)

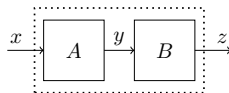


- Feedback composition

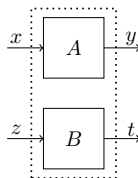


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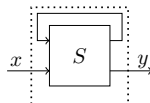
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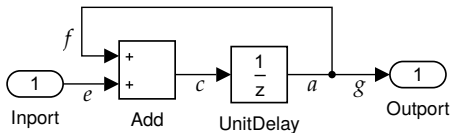


- Feedback composition



How to define feedback composition?

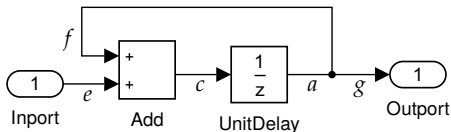
“Easy” feedback: “broken” by unit delays



$$a(k) = c(k - 1)$$

No instantaneous cyclic dependency.

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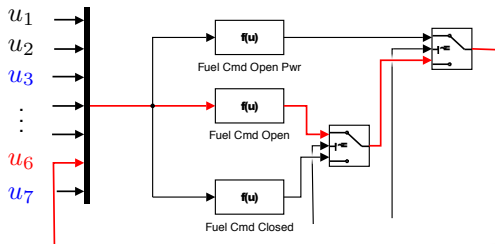
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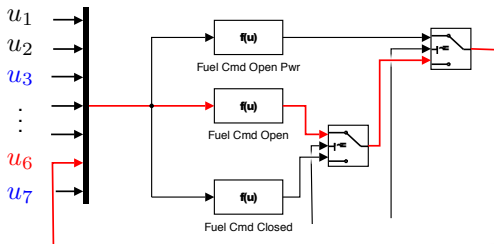
Handled in our earlier work on relational interfaces:

Tripakis, Lickly, Henzinger, Lee. *A Theory of Synchronous Relational Interfaces*, TOPLAS 2011.

“Not too hard” feedback



“Not too hard” feedback

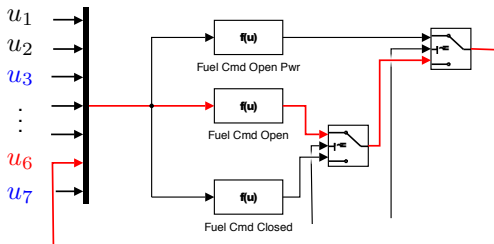


Output of FuelCmdOpen does not depend on u_6 :

$$\text{FuelCmdOpen}(u_1, u_2, \dots, u_7) = \frac{1}{14.7} (-0.366 + 0.08979u_7u_3 - 0.0337u_7u_3^2 + 0.0001u_7^2u_3)$$

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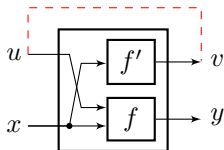
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Handled in recent work:

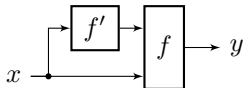
Dragomir, Preteasa, Tripakis. *Compositional Semantics and Analysis of Hierarchical Block Diagrams*, SPIN 2016.

“Not too hard” feedback: solution

If we know the block's **internals and input-output dependencies**:

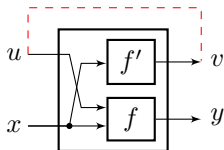


then this feedback reduces to serial composition:

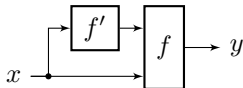


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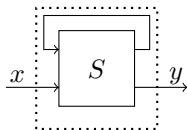


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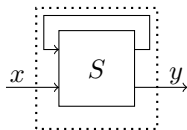
What if we don't know the internals of the block?

How to define general feedback?



How to define feedback for **any system** S ?

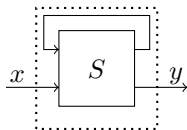
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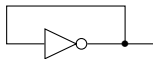
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A non-trivial problem.

Even for deterministic and input-receptive systems.

Feedback for **deterministic and input-receptive** systems

How to distinguish invalid feedbacks:



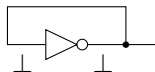
from valid ones:



Solution [Malik '94]: **fixpoint** semantics, starting with **unknown** values (“bottom” \perp).

Fixpoint semantics (also called *constructive semantics*)

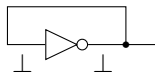
$$\neg \perp = \perp$$



Fixpoint stabilizes to \perp : feedback is invalid.

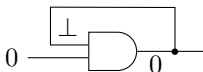
Fixpoint semantics (also called *constructive semantics*)

$$\neg \perp = \perp$$



Fixpoint stabilizes to \perp : feedback is invalid.

$$\perp \wedge 0 = 0$$



Fixpoint stabilizes to 0 : feedback is valid.

Limitations of constructive semantics

- It only applies to **deterministic and input-receptive** systems (i.e., total functions)
We also want to handle non-deterministic and non-input-receptive systems (partial relations).

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We also want to handle non-deterministic and non-input-receptive systems (partial relations).
- There is **no refinement** in existing constructive semantics frameworks.
*We not only have refinement, we also want refinement to be **preserved** by feedback.*

Preservation of refinement by composition

We want feedback to be **compositional**, i.e., for a composition operator \circ :

If $A \sqsubseteq A'$ (A' refines A) and $B \sqsubseteq B'$ then $A \circ B \sqsubseteq A' \circ B'$.

This property is necessary for substitutability.

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In particular, for feedback, we want:

If $A \sqsubseteq A'$ then $\text{feedback}(A) \sqsubseteq \text{feedback}(A')$.

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This is not easy to achieve – usual definitions don't work [TOPLAS 2011, FPS 2014, LICS 2016].

Results of LICS 2016 paper

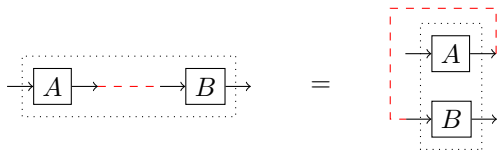
- 1 Two feedback operators:
 - **Instantaneous feedback**: applies to stateless (“memoryless”) systems.
 - **Feedback with unit-delay**: applies to stateful systems; instantaneous dependencies “broken” by a unit delay.
- 2 Both operators can handle **non-deterministic and non-input-receptive** systems.
- 3 Both operators proven to be compositional: **they preserve refinement**.

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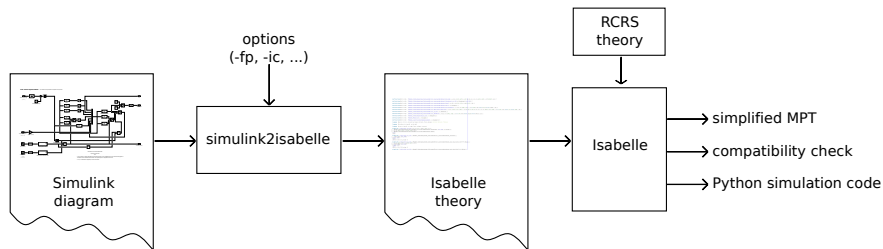
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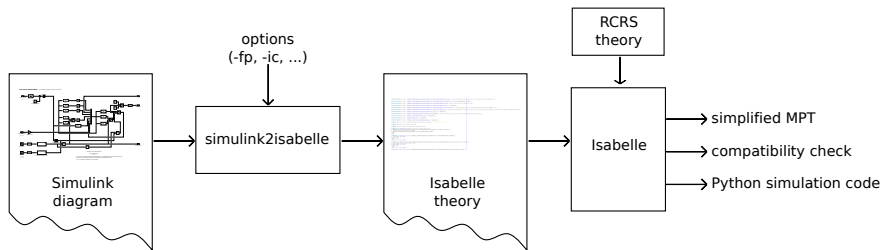
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- 5 Serial composition = parallel composition followed by feedback.



Toolset – downloadable from: rcrs.cs.aalto.fi



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Another non-trivial problem: *Translation of arbitrary block diagrams to terms using the three primitive composition operators (serial, parallel, feedback).*

See: Dragomir, Preteasa, Tripakis. *Compositional Semantics and Analysis of Hierarchical Block Diagrams*, SPIN 2016.

Thank you – questions?

Bibliography:



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